### A *Formal Methods* Approach Towards Trustworthy Global Computing

#### Diego Latella

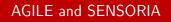
Consiglio Nazionale delle Ricerche Ist. di Scienza e Tecnologie dell'Informazione "A. Faedo" Formal Methods && Tools Lab

Bridging Global Computing with Grid November 28-29, 2006 Many concepts and ideas presented here are based on results on formal modeling and analysis of stochastic aspects of system behaviour achieved in the last few years by many colleagues and friends, a.o.:

- C. Baier et al. (Tec. Univ. of Dresden, D);
- M. Bernardo et al. (Univ. of Urbino, IT);
- R. Gorrieri et al. (Univ. of Bologna, IT);
- B. Haverkort et al. (Univ. of Twente, NL);
- H. Hermanns (Univ. of Saarbruecken, D);
- J. Hillston et al. (Univ. of Edinburg, UK);
- J.P. Katoen et al. (Univ. of Aachen, D);
- M. Kwiatkowska et al. (Univ. of Birmingham, UK);
- ... and many others!

### AGILE and SENSORIA

The focus on GC and SOC is the subject of cooperative work in the context of the EU Projects



Many thanks to:

- R. De Nicola (Univ. Firenze)
- J. P. Katoen (Univ. Aachen)
- M. Loreti (Univ. Firenze)
- M. Massink (CNR/ISTI, Pisa)

#### Background;

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#### Ochallenges;

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- Background;
- Ochallenges;
- (Approaches to) Solutions and Opportunities;

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- Open Issues;

"All engineering disciplines make progress by employing mathematically based notations and methods. Research on 'formal methods' follows this model and attempts to identify and develop mathematical approaches that can contribute to the task of creating computer systems"

[C. Jones 2000]

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Attempt to provide the (software) engineer with "concepts and techniques as thinking tools, which are clean, adequate, and convenient, to support him (or her) in describing, reasoning about, and constructing complex software and hardware systems"

[W. Thomas 2000]

Applying 
$$\left\{ \begin{array}{c} \text{Logic in} \\ \hline \text{Theoretical} \end{array} \right\}$$
 Computer Science

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- Automatic, often push-botton, Software Tool Support

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• classical issues like completeness.

Model-checking models of *trillions* of states (or more ... 10<sup>30</sup>), e.g.

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Complex control software for space applications (e.g. NASA Mars rovers [COMPUTER, Jan. 04])

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> Bill Gates, April 18, 2002. Keynote address at WinHEC 2002

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based on the solid framework of Mathematical Logic and equipped with efficient decision procedures (e.g. stochastic model-checkers)

D. Latella (CNR/ISTI)

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• Parallelism

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- Distribution

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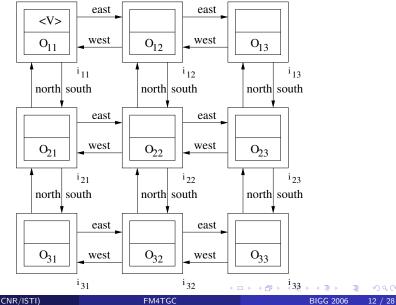
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 Resources which can be down-loaded using patter-matching mechanisms (discovery, SLA);

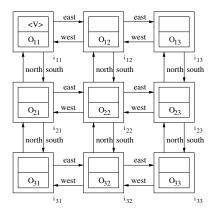
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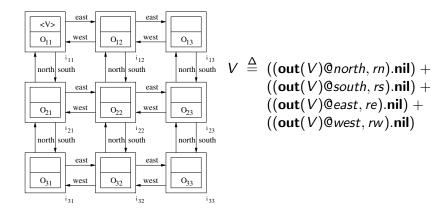
- Resources which can be down-loaded using patter-matching mechanisms (discovery, SLA);
- Uncertainty on operation execution times and/or possible choices by means of random variables and weights/probabilities (partial/approx. knowledge, performance/dependability analysis)



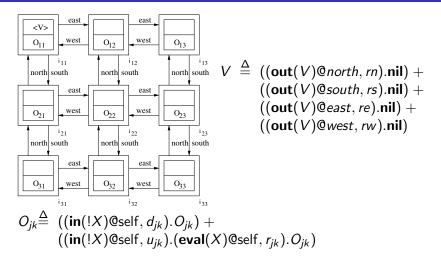
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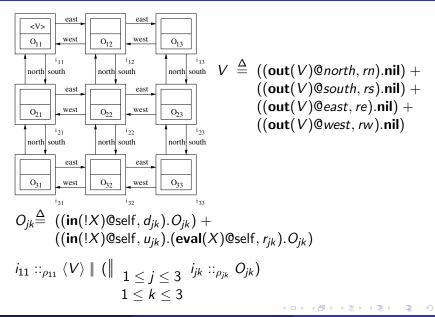


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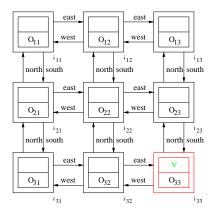


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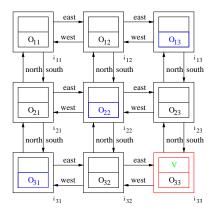




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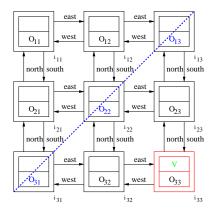
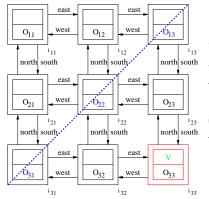


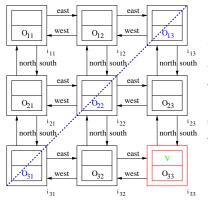
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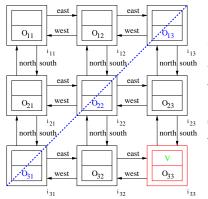


The probability is less then 0.2 that the infection develops (i.e. the virus is running) at site  $i_{33}$  by time *t*, assuming that at time 0 the only site infected is (i.e. the virus is stored at)  $i_{11}$ .



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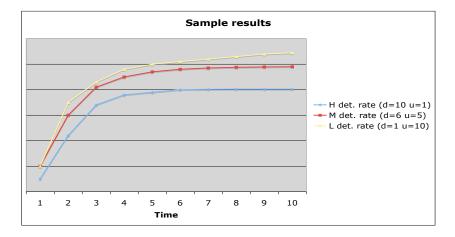
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... and by means of Automatic Stochastic Model-checking ....

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### Challenges, Solutions & Opportunities: Automatic Analysis



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Logical characterization (and automatic verification) of Steady State Probability measures

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In the long term,



Logical characterization (and automatic verification) of Steady State Probability measures

In the long term, the probability is greater than 0.9

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In the long term, the probability is greater than 0.9 that value f is present at site i.

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#### $S_{<0.2}(P@i)$

P could be a malicious (or faulty) process: the formula would characterize the average fraction of time the infection (or faulty component) is active at site i.

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Transient probability measures

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The probability is

#### $\mathcal{P}$ ( )

Transient probability measures

The probability is at least 0.8

 $P_{>0.8}($  )

Transient probability measures

The probability is at least 0.8 that the system is not down

 $\mathcal{P}_{>0.8}( \neg down)$ 

Transient probability measures

The probability is at least 0.8 that the system is not down at time t

 $\mathcal{P}_{>0.8}(\diamond^t \neg down)$ 

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The probability is less than 0.01

$$P_{<0.01}($$
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More general path-based measures

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Transient probability measures

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 $\mathcal{P}_{>0.8}(\diamond^t \neg down)$ 

Instantaneous availability at time t.

More general path-based measures

The probability is less than 0.01 that the system goes down within time t without having first raised an alarm signal

 $\mathcal{P}_{<0.01}(\neg alarm \ \mathcal{U}^{< t} \ down)$ 

**Nested** measures

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In equilibrium, the probability is

## *S* ( )

Nested measures

In equilibrium, the probability is at least 0.87,

## S<sub>>0.87</sub>( )

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In equilibrium, the probability is at least 0.87, that in at least 75% of the cases

$$S_{>0.87}(P_{>0.75}())$$

Nested measures

In equilibrium, the probability is at least 0.87, that in at least 75% of the cases the system will have resumed

 $S_{>0.87}(P_{>0.75}( \neg down))$ 

Nested measures

In equilibrium, the probability is at least 0.87, that in at least 75% of the cases the system will have resumed within 500 time units

 $S_{>0.87}(P_{>0.75}(\textit{down } U^{<500} \neg \textit{down}))$ 

Nested measures

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CTL as a special case

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Functional and non functional issues of behaviour integrated in the same description language

D. Latella (	CNR/ISTI)
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• For example

• For example stochastic simulators

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for state-transition models specified for instance in stochastic process-algebra(-like) languages

• For example stochastic simulators and (stochastic) mode-checkers

for state-transition models specified for instance in (stochastic) process-algebra(-like) languages and requirements characterized via (stochastic) temporal logics.

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Translations from e.g. UML, BPEL4WS, GRID-Skeletons to such languages and logics;

[e.g. EU-HIDE, EU-SENSORIA]

# Challenges, Solutions & Opportunities: Automatic Tool Support

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• Link to implementation:

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# Challenges, Solutions & Opportunities: Automatic Tool Support

• For example stochastic simulators and (stochastic) mode-checkers

for state-transition models specified for instance in (stochastic) process-algebra(-like) languages and requirements characterized via (stochastic) temporal logics.

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• Novel techniques and technologies are being developed

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e.g. [Kwiatkowska et al. 05]

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• Spatial, stochastic logics for mobility

[Theme of the 2nd National Software Summit hosted by the Center for National Software Studies Washington D.C. - May 10-12, 2004 ]

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